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COMPUTERIZED AUDIO BROADCASTING FRAMEWORK UTILIZING CODED OFDM

*SHAIK SHEHBAAZ BANU ** K. RAJYALAKSHMI

M.TECH Student, Department of ECE QUBA COLLEGE OF ENGINEERING AND TECHNOLOGY

Associate Professor QUBA COLLEGE OF ENGINEERING AND TECHNOLOGY

Abstract

In this paper we exhibit an investigation of BitError Rate (BER), for Digital Audio Broadcasting(DAB) framework, utilizing Coded OFDM withvarious channel coding plans. Examination is donefor convolutional coded and turbo codedinformation in an Additive White GaussianChannel (AWGN) in light of various imperativelengths and code generator polynomials utilized forcoding. A relative study on the computationalmultifaceted nature is likewise done by applying ansound flag and measuring the informationpreparing time per outline, on PCs with variousprocessor speeds. It is demonstrated that a codingincrease of roughly 6 dB is accomplished utilizingturbo coding when contrasted with convolutioncoding, at an expense of higher computationalmultifaceted nature.

KEYWORDS– DAB, OFDM, Convolutional Codes,Turbo Codes.

1. INTRODUCTION

The prerequisite of versatility whileassociated with system is energizing the developmentof remote correspondence. The routine simpletransmission systems don't perform well in versatileenvironment, since appropriate procedures tomoderate the impacts of multipath spread instigatedframeworks. Orthogonal Frequency DivisionMultiplexing (OFDM) is one such strategy to battlethe impact of multipath blurring, recurrence specific

blurring and Intersymbol Obstruction (ISI) [1].OFDM Diminishes the measure of equipmentexecution since multiplexing and sifting operationscan be performed by utilizing the Fast FourierTransform (FFT). This wipes out the need numerousoscillators at the transmitter and synchronizingcircles at the beneficiary. Because of the cyclic
augmentation of sign period into a gatekeeperinterim, OFDM framework is reasonable for SingleFrequency Networks (SFN) [5].

In this paper an OFDM application standardcalled Computerized Audio Broadcasting (DAB)framework model is executed in Matlab/Simulinkenvironment. The execution of this framework over achannel irritated by AWGN clamor is considered.Coded Orthogonal Frequency Division Multiplexing(COFDM) procedure is considered in whichconvolutional codes and turbo codes are utilizedwhat's more, figured the subsequent piece mistakerates (BER). The variety in BER is dissected takinginto account diverse coding parameters. A sound signis transmitted and information preparing time peredge is measured and thought about for diversechannel coding plans.

II. SYSTEM MODEL OF DAB USINGCODED OFDM

A. A Simplified DAB Block Diagram

A general block diagram of the DigitalAudio Broadcasting transmission system is shown inFig. 1. The simple sign is encoded and connected tochannel encoder. After channel coding the bit streamsare QPSK mapped. The information is then gone toOFDM generator. The high information rate bitstream is separated into "N" parallel informationsurges of low information rate and separatelyregulated on to orthogonal subcarriers which isacknowledged utilizing IFFT calculation.Orthogonality of the subcarriers accomplishes zeroInter Image Interference, hypothetically [1]. At longlast, the OFDM image is given cyclic prefix and thefinished Touch outline structure is transmittedthrough an AWGN channel.

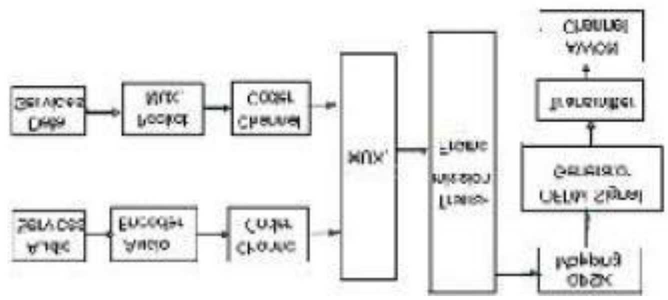

Figure 1. DAB transmitter – Block Diagram.

B. DAB Transmission Modes

DAB system has four transmissionmodes, each with its own set of parameters,shown in Table-I [12]. In this paperTransmission Mode-I is selected for simulation.

TABLE I. Dab Transmission Modes

Trans-mission Mode	No. of Sub-carriers	Sub carrier spacing	FFT Length	Maximum Radio Frequency
TM I	1536	1 KHz	2048	≈ 375 MHz
TM II	384	4 KHz	512	≈ 1.5 GHz
TM III	192	8 KHz	256	≈ 3 GHz
TM IV	768	2 KHz	1024	≈ 750 MHz

III. CHANNEL CODING

A. Convolutional Encoding & ViterbiDecoding

A convolutional encoder comprises of aM-stage shift register with "k" inputs, endorsedassociations with "n" modulo-2 adders andmultiplexer that serializes the yields of theadders. Here the encoder chose has k=1, ie; theinfo succession touches base on a solitaryinformation line. Subsequently the code rate isgiven by r = 1/n. In an encoder with a M stageshift enroll, the memory of the coder approachesM message bits and K = (M+1) movements arerequired before a message bit that has enteredthe movement register can at long last exit. Thisparameter K is alluded to as the imperativelength of the encoder.

The channel coding utilized for standardDAB comprises of code rate ½, memory 6,convolutional code with code generatorpolynomials 133 and 171 in octal organization[2]. For Spot lower code rates give betterexecution. Subsequently in this work, encoderwith code rate=w̅ is chosen. One suchconvolutional encoder is appeared in Fig. 2. Thequantity of registers=6. Consequently theimperative length K=7. Generator Polynomialsare 171, 133 and 115 in octal organization.Reproduction is done for different estimations oflimitation length and generator polynomials,which are given in Table-III.

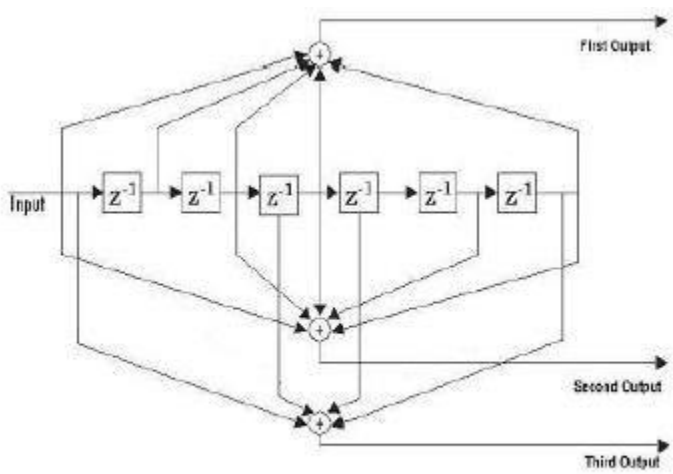

Figure 2. A rate w̅ convolutional encoder with constraint length, K=7.

B. Parallel Concatenated ConvolutionalTurbo Coding & Decoding

Parallel Concatenated Convolutional turbocode (PCC turbo code) comprises of two or moreRecursive Systematic Convolutional (RSC) codersworking in parallel [8]. The reason for interleaver isto offer each encoder an irregular variant of the databringing about equality bits from each RSC that arefree.

On the receiving side there are same numberof decoders as on the encoder side, each working onthe same information and an independent set of paritybits.

In this work, to give same code rate to turboencoder as on account of convolutional encoder, aparallel connection of two indistinguishable RSCencoders are utilized which gives a code rate of w̄.One such turbo encoder is appeared in Fig. 3, wherethe quantity of registers in each RSC encoder=2.

Henceforth the limitation length K=3. Generatorpolynomials are 7 and 5 in octal arrangement. Thenumber 7 indicates the input polynomial. ʌ is thearbitrary interleaver. Reenactment is completed fordifferent estimations of requirement length, generatorpolynomials and input polynomials, which are givenin Table-III.

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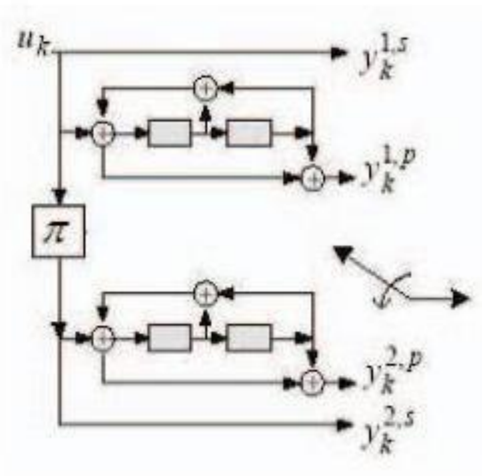

Figure 3. A rate w̄ turbo encoder with 2 parallel recursive systematic convolutional encoders, each with constraint length, K=3.

The inputs are data bits and called uk. Theyields are code bits. Of these, the yield of firstencoder, yk1, s is known as the precise piece, and itis the same as the information bit. The second yieldbit, yk1, p is the principal equality bit which isrecursive deliberate piece. An interleaver, indicatedby ʌ, is put in the middle of the two encoders toguarantee that the information got by the secondencoder is measurably autonomous. The third yieldbit, yk2, p is the second equality bit which isadditionally a recursive deliberate piece. The fourthyield yk2, s is deterministically reshuffling adaptationof yk1, s, which is not transmitted.

For decoding, the Viterbi Algorithm is notsuited to generate the A-Posteriori-Probability (APP)or soft decision output for each decoded bit. HereMaximum-A-Posteriori (MAP) algorithm is used forcomputing the metrics. Block diagram of turbodecoder is shown in Fig. 4.

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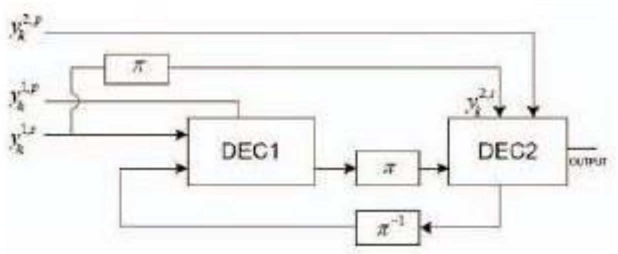

Figure 4. Turbo decoder – Block Diagram.

In Fig. 4, DEC1 and DEC2 are 2 APPdecoders. ʌ and ʌ -1 are random interleaver anddeinterleaver respectively [14]. The symbol vectorsent for each time are described by yk=(yk1, s, yk1,p, yk2, p). The goal is to take these and make a guessabout the transmitted vector and hence code bitswhich in turn decode uk, the information bit.

IV. SIMULATION MODEL

A. Simulation Parameters

The simulation parameters are shown in Table-II [1].The different channel coding schemes and itsparameters used for the analysis are given in Table-III. Even though, a complete DAB system consists ofa multiplex of many information service channels,here, for the purpose of analysis, only a single audiosignal is selected for transmission**.**

TABLE II. Simulation Parameters

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Transmission Mode	Mode I
No. of sub-carriers	1536
Transmission frame duration (F_t)	96 ms.
OFDM Symbols per Transmission Frame	76
Sample Time (T_s)	0.48828 µs
Frame length	196608 (or F_t/T_s)
FFT length	2048
Guard interval (Cyclic Prefix)	504
OFDM length	2552
Channel Coding schemes Used	Convolutional coding (rate ⅓), Turbo coding (rate ⅓)
Modulation	QPSK
Channel	AWGN

TABLE III. Channel Coding Parameters

Channel Coding types	Constraint length	Code generator polynomials (Octal format)	Feedback Polynomial
Convolutional Coding	3	7, 6, 5	--
	4	15, 13, 11	--
	5	34, 27, 23	--
	6	71, 57, 47	--
	7	171, 133, 115	--
Turbo Coding	3	7, 5	7
	4	15, 13	15
	5	34, 27	34
	6	71, 57	71
	7	171, 133	171

B. Simulation Block Diagram

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Fig. 5 and Fig. 6 show the simulation models forDAB using convolutional coding and turbo codingrespectively. Simulations are carried out usingMatlab/Simulink.

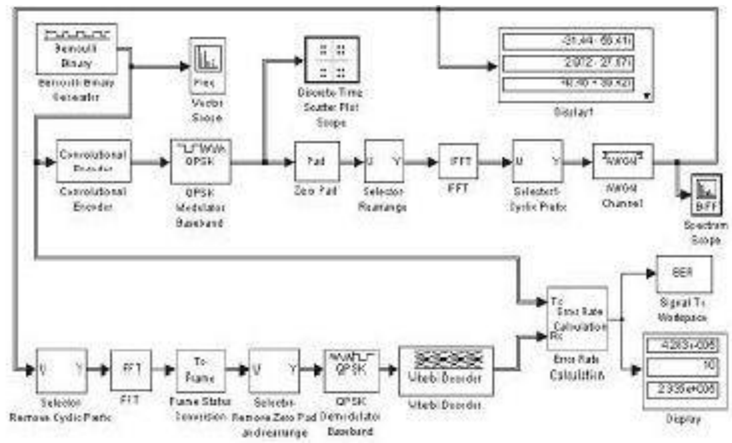

Figure 5. Simulation model for DAB transceiver using convolutional coding & viterbi decoding.

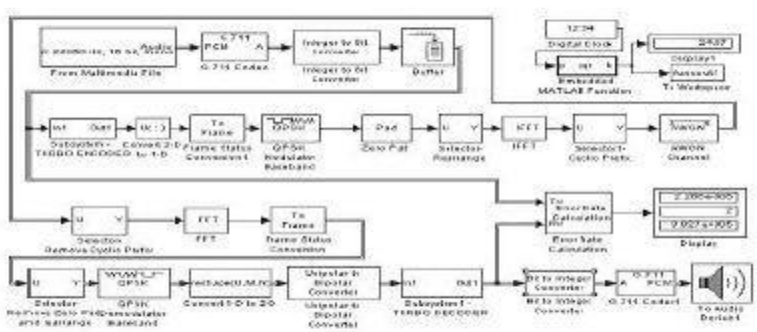

Figure 6. Simulation model for DAB transceiver using turbo coder and decoder. An audio signal is inputted and analyzed.

V. RESULTS

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Coding Scheme	Frame Processing time on low speed computer (milli seconds)	Frame Processing time on high speed computer (milli seconds)
Conv_CnstrLn3	3.93	2.59
Conv_CnstrLn4	4.18	2.81
Conv_CnstrLn5	4.60	3.13
Conv_CnstrLn6	5.51	3.86
Conv_CnstrLn7	7.07	5.18
Turbo_CnstrLn3	28.36	16.79
Turbo_CnstrLn4	44.43	27.66
Turbo_CnstrLn5	77.89	50.20
Turbo_CnstrLn6	142.50	92.57
Turbo_CnstrLn7	270.90	180.30

TABLE IV. Frame Processing Time Comparison

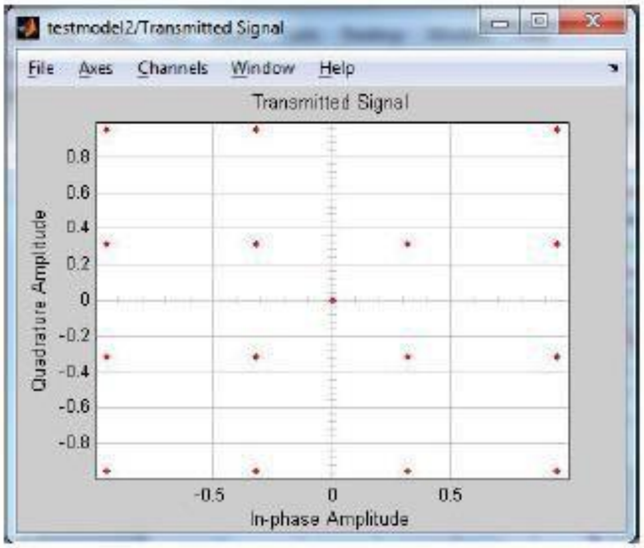

Figure 7. Transmitted Signal

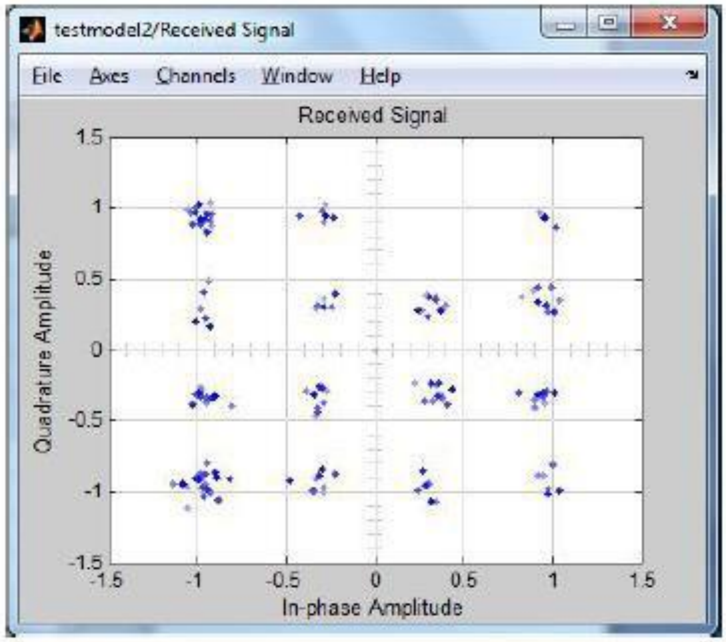

Figure 8. Received Signal

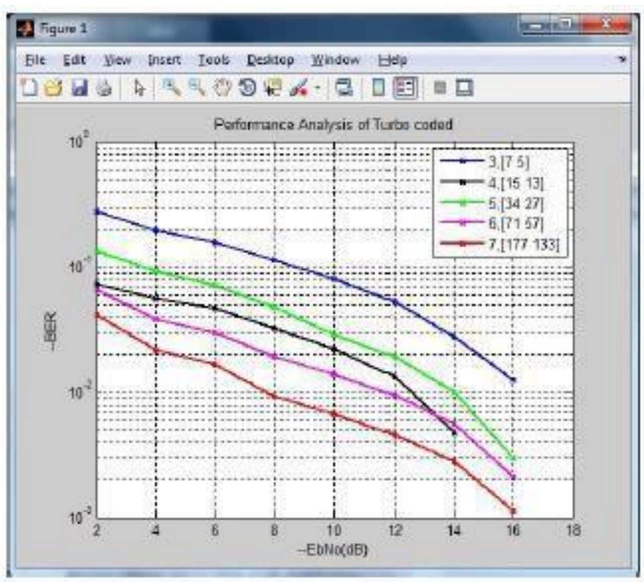

Figure 9. Performance Analysis Of Turbo Codes

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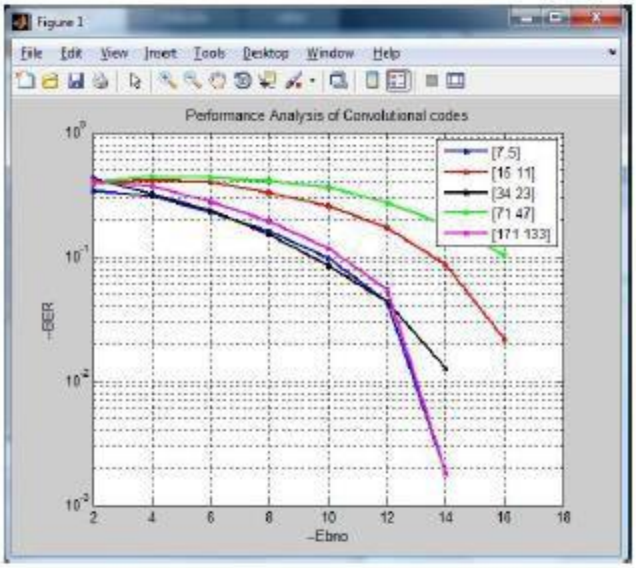

Figure10. Performance Analysis Of Convolutional Codes

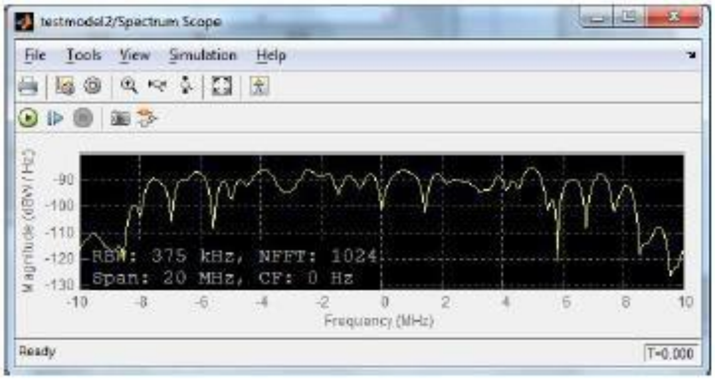

Figure 11. Frequency spectrum

VI. CONCLUSION

Computerized Audio Broadcasting frameworkutilizing Coded OFDM is executed and contemplatedover an AWGN channel. Bit-Error-Rate (BER) ismeasured and looked at by utilizing mistake adjustingcodes like, Convolutional Code and ParallelConcatenated Convolutional Turbo Code. A decentBER for sound is thought to be 10-4. Utilizing turbocoding, it is almost accomplished with an Eb/No of 3dB. A coding increase of about 6 dB is accomplishedutilizing turbo coding, when contrasted withconvolutional coding, at an expense of highcomputational many-sided quality. Likewise,reenactment is done on low speed and rapid PCs andedge preparing time is measured as a section of studyon Quality of Service (QoS). It is demonstrated that,slightest complex turbo code obliges 3 to 4 times thepreparing time taken by most astounding complexconvolutional code. In this way coding increase isaccomplished utilizing turbo code by trading off oncomputational time required.The burdens of the customary codes like convolutional codes is that, with an end goal toapproach the hypothetical point of confinement forShannon's channel limit, we require to expand theimperative length of a convolutional code, which,thus, causes the computational intricacy of a mostextreme probability decoder to incrementsexponentially. At last we achieve a point whereunpredictability of the decoder is high to the pointthat it gets to be hard to figure it out physically
furthermore there is no impressive diminishment inBER. Further coding addition is conceivable withturbo codes with sensible coding and decipheringmultifaceted nature. In this paper, it is intended tolegitimize these conclusions with reenactments.

The channel chose just presents Gaussian clamor. Inany case, issues confronted by blurring and multipathcan be examined by picking other channel models.Rather than QPSK adjustment conspire, the sameframework can be examined utilizing other tweakprocedures like DQPSK also, QAM. Rather thanparallel linked convolutional turbo codes, serialconnected convolutional turbo code additionally canbe executed and examined.

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